

3C03 Concurrency: Liveness & Progress

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Outline

- Liveness
- Progress
- Progress Specification in FSP
- Progress-Analysis of LTS
- Priorities

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Motivation

- Problem with single lane bridge:
- Cars cannot pass from north to south if there is a continuous stream of cars from south to north!
- We would like to guarantee that cars will eventually cross the bridge.
- In more general terms this is referred to as liveness

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Liveness

- A <u>liveness property</u> asserts that something good eventually happens.
- We want to specify liveness for our FSP models
- We want to analyze our FSP models to be certain that the liveness properties hold
- General form of liveness requires consideration of temporal precedence relationship between states
- We use more restricted form of progress

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Progress

- A <u>progress property</u> asserts that whatever state a system is in, it is always the case that a specified action will eventually be executed
- Progress is the opposite form of starvation
- Notion of progress is sufficiently powerful to capture wide range of liveness properties
- Progress properties are simple to specify in FSP

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Progress Properties in FSP

- Specification of progress needs assumption of a fair scheduling policy.
- If a transition from a set is chosen infinitely often and every transition in the set will be executed infinitely often, the scheduling policy is said to be <u>fair</u>.
- progress P={a1,a2,...,an} defines a progress property P which asserts that in an infinite execution at least one of the actions a1, a2, ..., an will be executed infinitely often.

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Example: Tossing Coins

```
COIN = ( toss -> heads -> COIN | toss -> tails -> COIN).

toss
toss
toss
tails

progress HEADS = {heads}
```

progress HEADS = {heads}
progress TAILS = {tails}

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Example: Tossing Trick Coins

```
TWOCOIN = (pick->COIN | pick->TRICK),
         = ( toss -> heads -> COIN
COIN
              toss -> tails -> COIN),
         = (toss->heads->TRICK).
progress HEADS = {heads}
progress TAILS = {tails}
             pick
                            toss
             toss
                             toss
     pick
             heads
                             heads
                                  tails
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```



Progress Analysis

- We can automate analysis of progress properties
- A set of states where every state is reachable from every other state in the set and no state has transitions to states outside the set is a terminal set of states.
- Terminal set of states can be found using a graph algorithm that searches for a strongly connected component.

 LTSA

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Default Progress Properties

- <u>Default progress properties</u> assert in a system with fair choices that every action in the alphabet will be executed infinitely often.
- Default progress properties of example:

```
progress p1 = {pick}
progress p2 = {toss}
progress p3 = {heads}
progress p4 = {tail}
```

■ How many violations?

LTSA

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Priorities

- Default progress analysis of single lane bridge does not reveal violation. LTSA
- Problem is scheduling policy. Cars arriving in the south get 'priority' if there are already northbound cars on the bridge
- To detect such progress violations we have to reflect such priorities in the FSP model

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High Priority in FSP

■ ||C = (P||Q)<<{a1,...,an} specifies a composition in which the actions a1,...,an have higher priority than any other action in the alphabet of P||Q including the silent action tau. In any choice in this system which has one or more of the actions a1,...,an labelling a transition, the transitions labelled with lower priority actions are discarded.</p>

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Low Priority in FSP

■ ||C = (P||Q)>>{a1,...,an} specifies a composition in which the actions a1,...,an have lower priority than any other action in the alphabet of P||Q including the silent action tau. In any choice in this system which has one or more transitions not labelled by a1,...,an, the transitions labelled by a1,...,an are discarded.

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Simplification of LTS

- Priorities simplify the LTS resulting of the composition.
- Example:

```
NORMAL=(work->play->NORMAL)
| sleep->play->NORMAL).
| |HIGH=(NORMAL)<< {work}.
| LTSA</pre>
```

■ Use of priorities lead to more realistic liveness checks.

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Summary

- Liveness
- Progress
- Progress Specification in FSP
- Progress-Analysis of LTS
- **■** Priorities
- Next session: Progress Analysis of Single-Lane Bridge

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