

C340 Concurrency: Semaphores and Monitors

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Revised Lecture Plan

- 1 Introduction
- 2 Modelling Processes
- 3 Modelling Concurrency in FSP
- 4 FSP Tutorial
- 5 LTSA Lab
- 6 Programming in Java
- 7 Concurrency in Java
- 8 Lab: Java Thread Programming

- 9 Mutual Exclusion
- 10 Lab: Synchronization in Java
- 11 Semaphores and Monitors
- 12 Conditional Synchronization
- 13 Fairness & Liveness
- 14 Safety
- 15 Tutorial: Model Checking



Goals

Introduce concepts of

- Semaphores
- Monitors
- Conditional synchronisation
- Relationship to FSP guarded actions

Implementation in Java

- synchronised methods and private attributes
- single thread active in the monitor at any time
- wait, notify and notifyAll



Semaphores

Introduced by Dijkstra' in 1968 ADT with counter and waiting list *P/Wait/Down:* S/Signal/Up: if (threads wait) if (counter > 0) activate waiting counter-thread else else add caller to waiting list counter++



Semaphores and Mutual Exclusion

- One semaphore for each critical section
 Initialize semaphore to 1.
- Embed critical sections in wait/signal pair
- Example in Java:
 - Semaphore S=new Semaphore(1);
 - S.down();
 - <critical section>
 - S.up();



Evaluation of Semaphores

- + Nice and simple mechanism
- + Can be efficiently implemented



- Too low level of abstraction
- Unstructured use of signal and wait leads to spaghetti synchronisation
- Error prone and errors are dangerous
 - Omitting signal leads to deadlocks
 - Omitting wait leads to safety violations



Critical Regions

Guarantee mutual exclusion by definition

Note subtle difference to critical sections

language features implement critical regions

Example: Java synchronised method



Monitors

Hoare's response to Dijkstra's semaphores

- Higher-level
- Structured
- Monitors encapsulate data structures that are not externally accessible
- Mutual exclusive access to data structure enforced by compiler or language run-time





Monitors in Java

- All instance and class variables need to be private Or protected
 All methods need to be synchronised
 Example: semaphore implementation
- Use of Monitors: Carpark Problem



Carpark Problem

- Only admit cars if carpark is not full
 Cars can only leave if carpark is not empty
- Car arrival and departure are independent threads

Demo: CarPark



Carpark Model

Events or actions of interest:

- Arrive and depart
- Processes:
 - Arrivals, departures and carpark control

Process and Interaction structure:





Carpark FSP Specification

```
CARPARKCONTROL(N=4) = SPACES[N],
SPACES[i:0...N] =
       (when(i>0) arrive-> SPACES[i-1]
       when(i<N) depart-> SPACES[i+1]
       ARRIVALS = (arrive-> ARRIVALS).
DEPARTURES = (depart -> DEPARTURES).
  CARPARK =
 (ARRIVALS | | CARPARKCONTROL | | DEPARTURES ).
```

12



Java Class Carpark

public class Carpark extends Applet final static int N=4; public void init() { CarParkControl cpk = new CarParkControl(N); Thread arrival, departures; arrivals=new Thread(new Arrivals(cpk)); departures=new Thread(new Departures(cpk)); arrivals.start(); departures.start();



Java Classes Arrivals & Departures

```
public class Arrivals implements Runnable {
CarParkControl carpark;
Arrivals(CarParkControl c) \{carpark = c;\}
public void run() {
 while (true) carpark.arrive();
class Departures implements Runnable {
public void run() {
 while (true) carpark.depart();
```



Java Class CarParkControl (Monitor)

```
class CarParkControl {// synchronisation?
 private int spaces;
 private int N;
 CarParkControl(int capacity) {
    N = capacity;
    spaces = capacity;
  synchronized public void arrive() {
    ... -- spaces; ... } {// Block if full?
  synchronized public void depart() {
    ... ++ spaces; ... {// Block if empty?
```



Problems with CarParkControl

- How do we send arrivals to sleep if car park is full?
- How do we awake it if space becomes available?
- Solution: Condition synchronisation





- Semaphores
- Monitors
- Next session:
 - Java condition synchronization
 - Relationship between FSP guarded actions and condition synchronization
 - Fairness and Starvation