



Concurrency Control



Motivation

- ***Components of distributed systems use shared resources concurrently:***
 - *Hardware Components*
 - *Operating system resources*
 - *Databases*
 - *Objects*
- ***Resources may have to be accessed in mutual exclusion.***



Motivation

- **Concurrent access and updates of resources may lead to:**
 - *lost updates*
 - *inconsistent analysis.*
- **Example for lost updates:**
 - *Cash withdrawal from ATM and concurrent*
 - *Credit of cheque.*
- **Example for inconsistent analysis:**
 - *Funds transfer between accounts of a customer*
 - *Sum of account balances (Report for Inland Revenue).*

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3



Motivating Examples

```
class Account {  
    protected:  
        float balance;  
    public:  
        float get_balance() {return balance;};  
        void debit(float amount){  
            float new=balance-amount;  
            balance=new;  
        };  
        void credit(float amount) {  
            float new=balance+amount;  
            balance=new;  
        };  
};
```

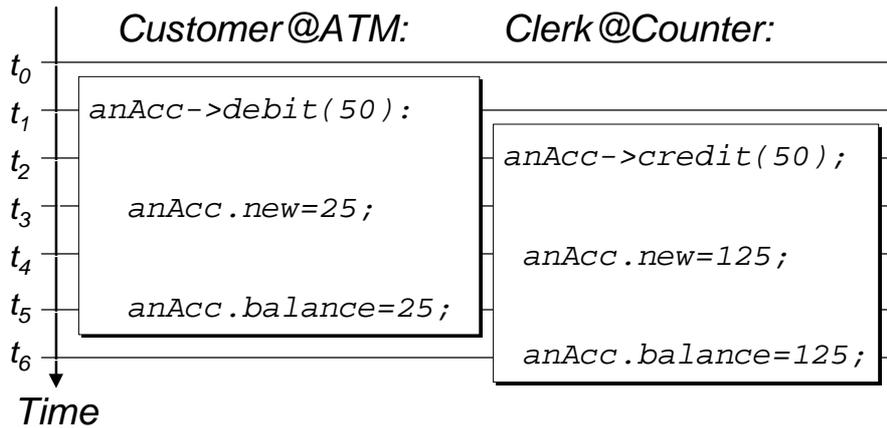
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4



Lost Updates

Balance of account *anAcc* at t_0 is 75



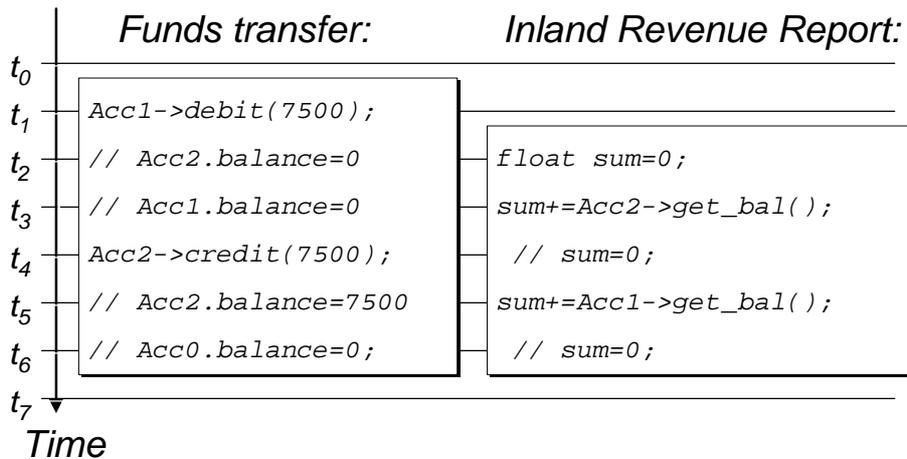
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5



Inconsistent Analysis

Balances at t_0 *Acc1*: 7500, *Acc2*: 0



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6



Two Phase Locking (2PL)

- **The most popular concurrency control technique. Used in:**
 - *RDBMSs (Oracle, Ingres, Sybase, DB/2, etc.)*
 - *ODBMSs (O2, ObjectStore, Versant, etc.)*
 - *Transaction Monitors (CICS, etc)*
- **Concurrent processes acquire locks on shared resources from lock manager.**
- **Lock manager grants lock if request does not conflict with already granted locks.**
- **Guarantees serialisability.**

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7



Locks

- **A lock is a token that indicates that a process accesses a resource in a particular mode.**
- **Minimal lock modes: read and write.**
- **Locks are used to indicate to concurrent processes the current use of that resource.**

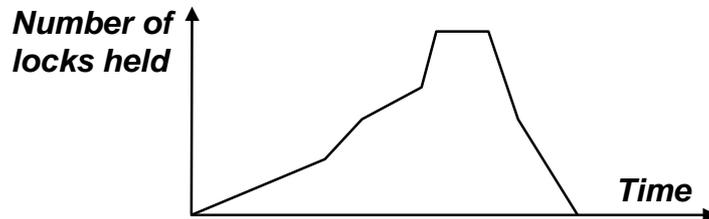
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8



Locking

- **Processes acquire locks before they access shared resources and release locks afterwards.**
- **2PL: Processes do not acquire locks once they have released a lock.**
- **Typical 2PL locking profile of a process:**



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9



Lock Compatibility

- **Lock manager grants locks.**
- **Grant depends on compatibility of acquisition request with modes of already granted locks.**
- **Compatibility defined in lock compatibility matrix.**
- **Minimal lock compatibility matrix:**

	Read	Write
Read	+	-
Write	-	-

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10



Locking Conflicts

- **Lock requests cannot be granted if incompatible locks are held by concurrent processes.**
- **This is referred to as a locking conflict.**
- **Approaches to handle conflicts:**
 - **Force requesting process to wait until conflicting locks are released.**
 - **Tell process or thread that lock cannot be granted.**

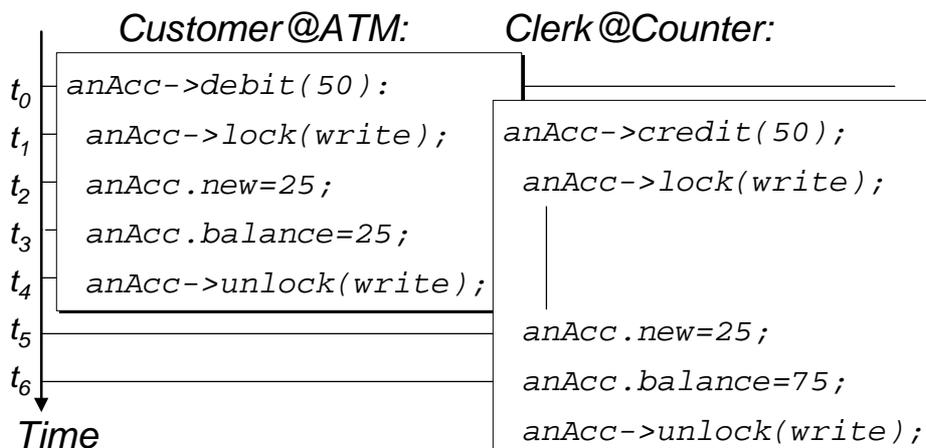
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11



Example (Avoiding Lost Updates)

Balance of account *anAcc* at t_0 is 75



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12



Deadlocks

- *2PL may lead to processes waiting for each other to release locks.*
- *These situations are called deadlocks.*
- *Deadlocks have to be detected by the lock manager.*
- *Deadlocks have to be resolved by aborting one or several of the processes involved.*
- *This requires to undo all the actions that these processes have done.*

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13



Locking Granularity

- *2PL applicable to resources of any granularity.*
- *High degree of concurrency with small locking granularity.*
- *For small granules large number of locks required.*
- *May involve significant locking overhead.*
- *Trade-off between degree of concurrency and locking overhead.*
- *Hierarchical locking as a compromise.*

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14



Hierarchical Locking

- **Used with container resources, e.g.**
 - *file (containing records)*
 - *set or sequence (containing objects)*
- **Lock modes intention read (IR) and intention write (IW).**
- **Lock compatibility:**

	IR	R	IW	W
IR	+	+	+	-
R	+	+	-	-
IW	+	-	+	-
W	-	-	-	-

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15



Transparency of Locking

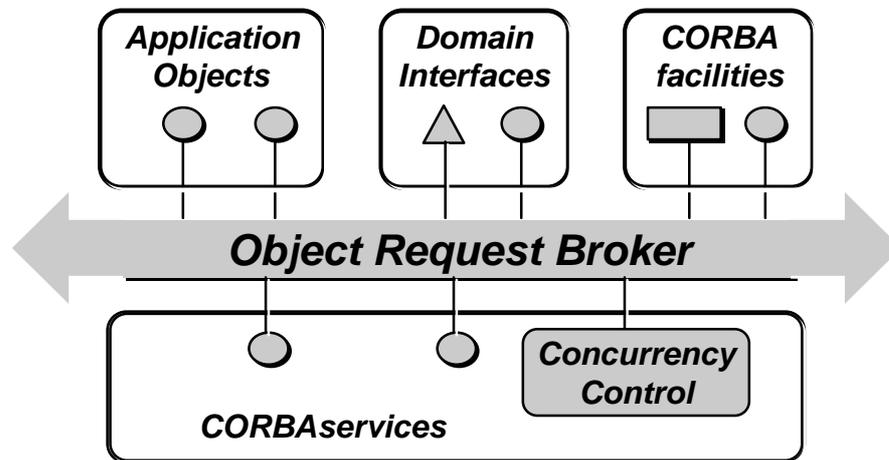
- **Who is acquiring locks?**
 - *Concurrency control infrastructure*
 - *Implementation of components*
 - *Clients of components*
- **First option desirable but not always possible:**
 - *Infrastructure must manage all resources*
 - *Infrastructure must know all resource accesses.*
- **Last option is undesirable and avoidable!**

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16



CORBA Concurrency Control Service



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Lock Compatibility Matrix

- **CORBA concurrency control service supports hierarchical locking.**
- **Upgrade locks for decreasing probability of deadlocks.**
- **Compatibility matrix:**

	IR	R	U	IW	W
IR	+	+	+	+	-
R	+	+	+	-	-
U	+	+	-	-	-
IW	+	-	-	+	-
W	-	-	-	-	-

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18



Locksets

- *A lockset is associated to a resource (usually in the implementation of that resource).*
- *Each shared resource has a lockset.*
- *Operations of that resource acquire locks before they access or modify the resource.*



The IDL Interfaces

```
interface LocksetFactory {
    LockSet create();
};

interface Lockset {
    void lock(in lock_mode mode);
    boolean try_lock(in lock_mode mode);
    void unlock(in lock_mode mode);
    void change_mode(in lock_mode held,
                    in lock_mode new);
};
```